

Recursion

Genome 559: Introduction to Statistical and
Computational Genomics

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The merge sort algorithm

1. *Split your list into two halves*
2. *Sort the first half*
3. *Sort the second half*
4. *Merge the two sorted halves, maintaining a sorted order*

Divide-and-conquer

- The basic idea behind the merge sort algorithm is to divide the original problem into two halves, **each being a smaller version of the original problem.**
- This approach is known as *divide and conquer*
 - Top-down technique
 - Divide the problem into **independent** smaller problems
 - Solve smaller problems
 - Combine smaller results into a larger result thereby “conquering” the original problem.

Merge sort – the nitty gritty

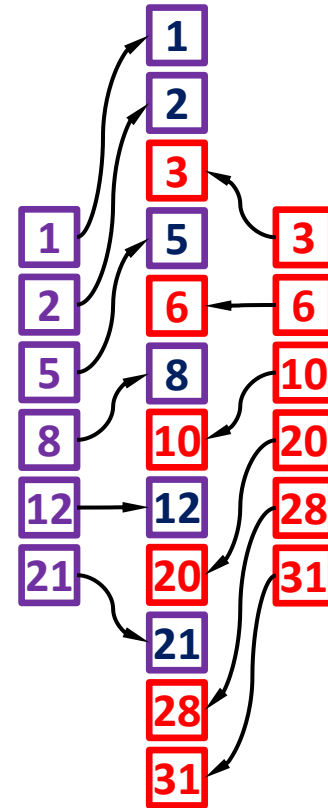
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That's simple

???

Careful bookkeeping, but still simple



If I knew how to sort, I wouldn't be here in the first place?!?

Merge sort – the nitty gritty

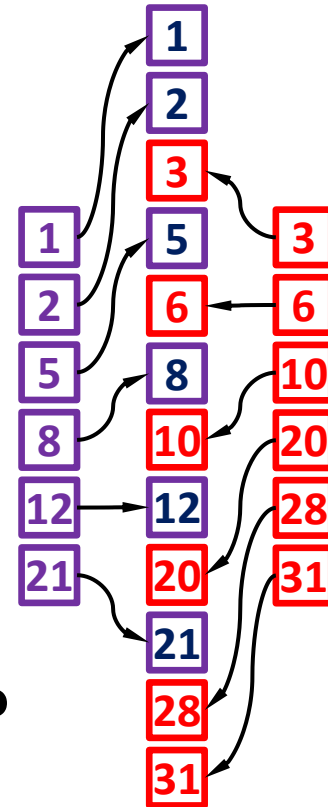
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So ...

how are we going to sort the two smaller lists?

**Here's a crazy idea:
let's use merge sort
to do this**

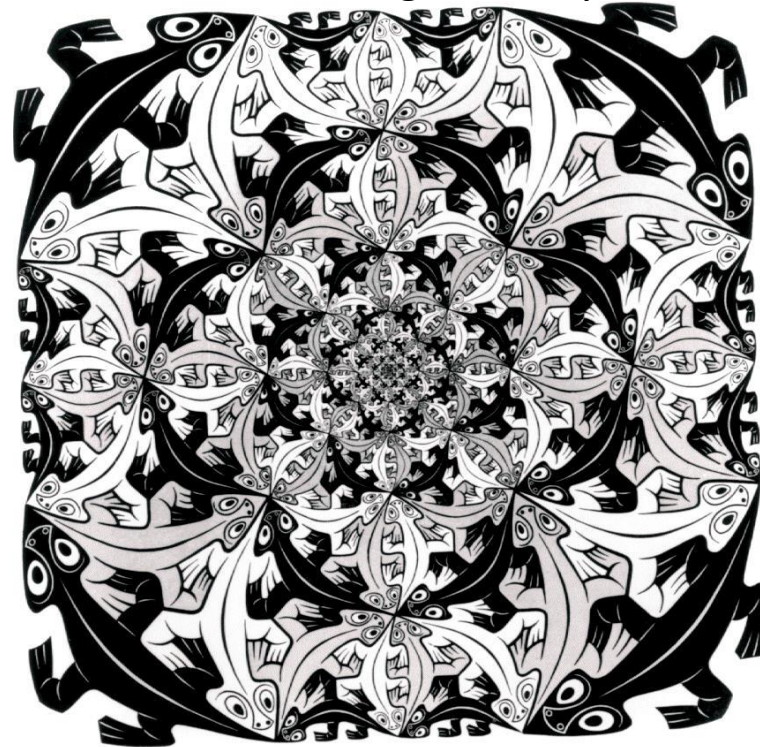
```
def mergeSort(list):  
    half1 ← first half of list  
    half2 ← second half of list  
    half1_sorted = mergeSort(half1)  
    half2_sorted = mergeSort(half2)  
    list_sorted = merge(half1_sorted, half2_sorted)  
    return list_sorted
```

You must be kidding, right?

- WHAT?
- This function has no loop?
- It seems to refer to itself!
- Where is the actual sort?
- What's going on???

```
def mergeSort(list):  
    half1 ← first half of list  
    half2 ← second half of list  
    half1_sorted = mergeSort(half1)  
    half2_sorted = mergeSort(half2)  
    list_sorted = merge(half1_sorted, half2_sorted)  
    return list_sorted
```

this is making me dizzy!



Let's take a step back ...

Factorial

- A simple function that calculates $n!$

```
# This function calculated n!  
def factorial(n):  
    f = 1  
    for i in range(1,n+1):  
        f *= i  
    return f
```

```
>>> print factorial(5)  
120  
>>> print factorial(12)  
479001600
```

- This code is based on the standard definition of factorial: $n! = \prod_{k=1}^n k$

Factorial

- But ... there is an alternative **recursive** definition:

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ (n-1)! \times n & \text{if } n > 0 \end{cases}$$

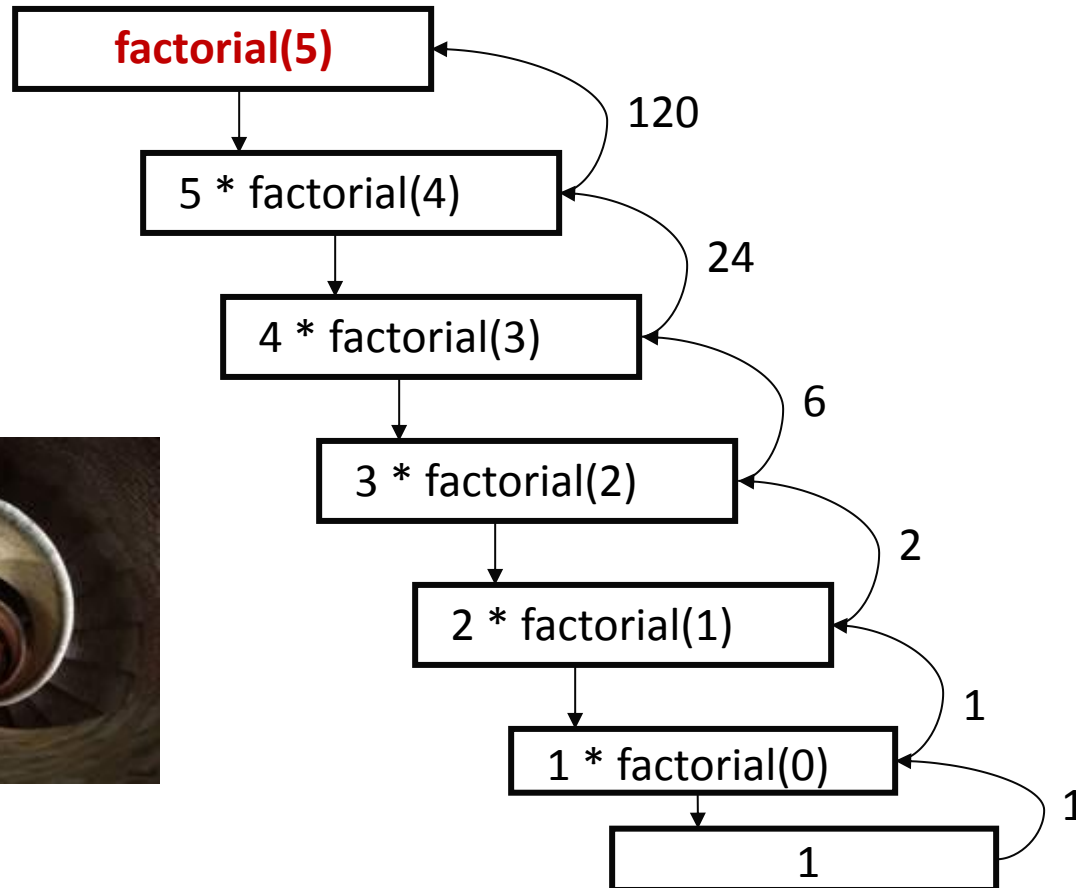
- So ... can we write a function that calculates n! using this approach?

```
# This function calculated n!  
def factorial(n):  
    if n==0:  
        return 1  
    else:  
        return n * factorial(n-1)
```

- Well ...
We can! It works! And it is called a ***recursive*** function!

Why is it working?

```
# This function calculated n!  
def factorial(n):  
    if n==0:  
        return 1  
    else:  
        return n * factorial(n-1)
```



Recursion and recursive functions

- **A function that calls itself**, is said to be a **recursive function** (and more generally, an algorithm that is defined in terms of itself is said to use recursion or be recursive)

(A call to the function “recurs” within the function; hence the term “recursion”)

- In many real-life problems, recursion provides an intuitive and natural way of thinking about a solution and can often lead to very elegant algorithms.

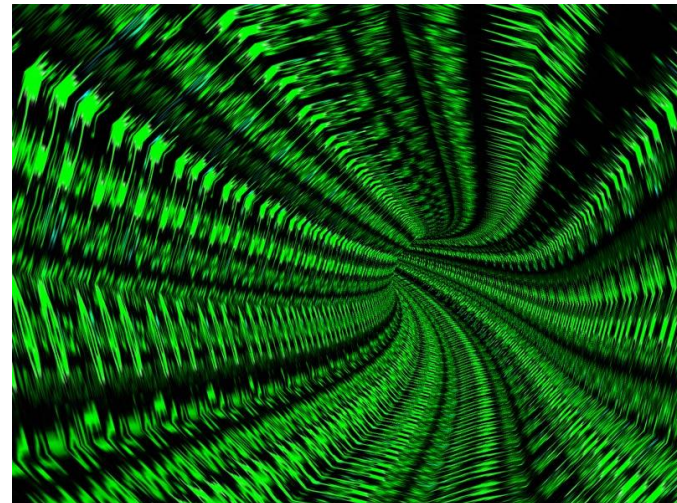
mmm...

- If a recursive function calls itself in order to solve the problem, isn't it circular?
(in other words, why doesn't this result in an infinite loop?)
- Factorial, for example, is not circular because we eventually get to $0!$, whose definition **does not rely** on the definition of another factorial and is simply 1.
 - This is called a **base case** for the recursion.
 - When the base case is encountered, we get a closed expression that can be directly computed.

Defining a recursion

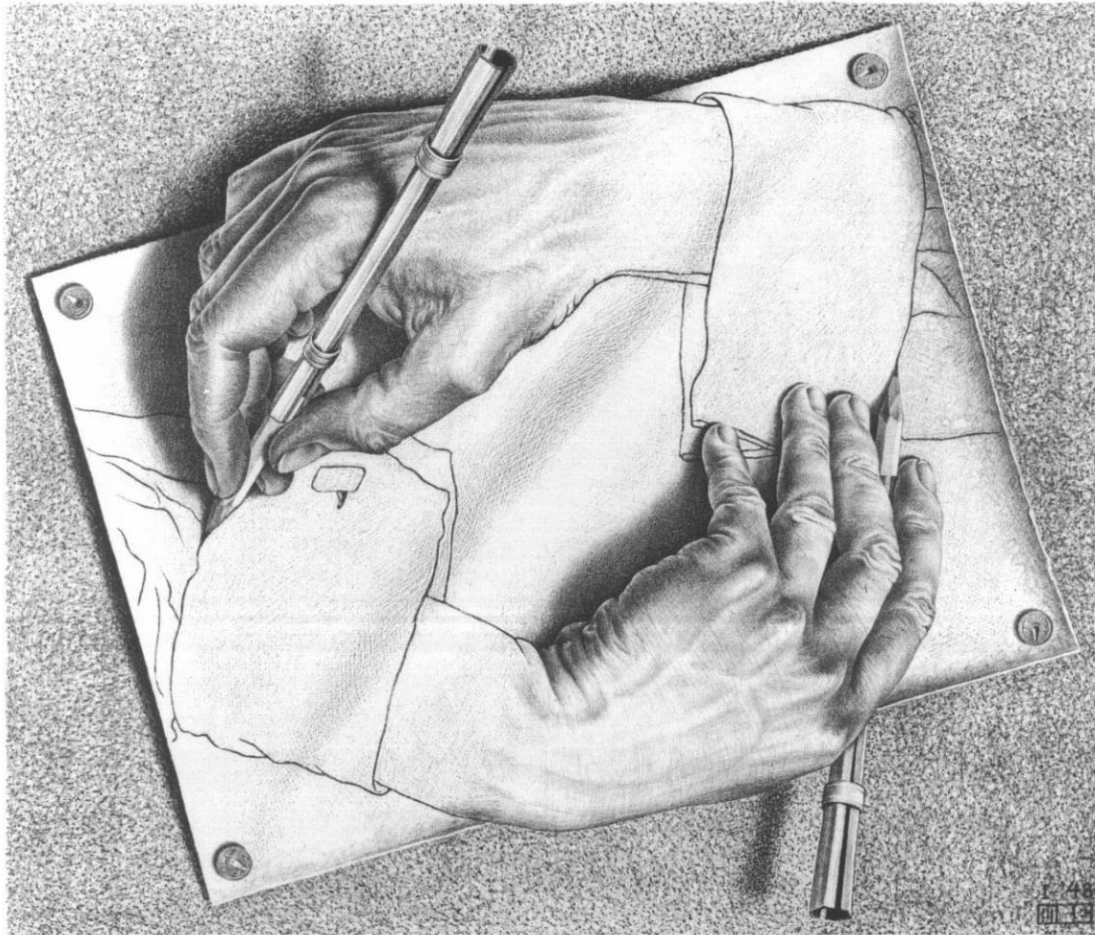
- Every recursive algorithm must have two key features:
 1. There are one or more **base cases** for which no recursion is applied.
 2. All recursion chains eventually end up at one of the base cases.

*The simplest way for these two conditions to occur is for each recursion to act on a **smaller** version of the original problem. A very small version of the original problem that can be solved without recursion then becomes the base case.*



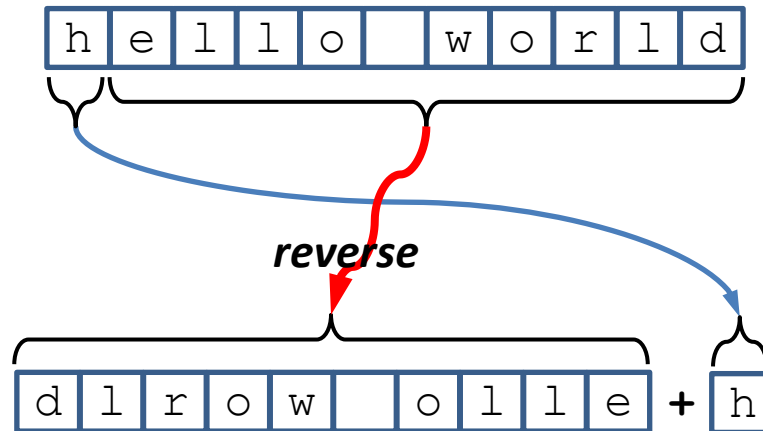
This is fun!

Let's try to solve (or at least think of)
other problems using recursion



String reversal

- Divide the string into first character and all the rest
- Reverse the “rest” and append the first character to the end of it



See how simple and elegant it is! No loops!

```
# This function reverses a string
def reverse(s):
    return reverse(s[1:]) + s[0]
```

*s[1:] returns all but the first character of the string. We **reverse** this part (s[1:]) and then concatenate the first character (s[0]) to the **end**.*

String reversal - D'oh!

```
# This function reverses a string
def reverse(s):
    return reverse(s[1:]) + s[0]
```

```
>>> print reverse("hello world")
Traceback (most recent call last):
  File "<stdin>", line 1, in <module>
  File "<stdin>", line 2, in reverse
  File "<stdin>", line 2, in reverse
  File "<stdin>", line 2, in reverse
  File "<stdin>", line 2, in reverse
  File "<stdin>", line 2, in reverse
  File "<stdin>", line 2, in reverse
  File "<stdin>", line 2, in reverse
  .
  .
  .
  File "<stdin>", line 2, in reverse
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  File "<stdin>", line 2, in reverse
  File "<stdin>", line 2, in reverse
RuntimeError: maximum recursion depth exceeded
```

What just happened? There are 1000 lines of errors!

String reversal – Duh!

- **Remember:** To build a correct recursive function, we need a **base case** that doesn't use recursion!

We forgot to include a base case, so our program is an infinite recursion. Each call to “reverse” contains another call to reverse, so none of them return.

Each time a function is called it takes some memory. Python stops it at 1000 calls, the default “maximum recursion depth.”

- **What should we use for our base case?**

String reversal - Yeah

- Since our algorithm is creating shorter and shorter strings, it will eventually reach a stage when S is of length 1 (one character).
- Since a string of length 1 is its own reverse, we can use it as the base case.

```
# This function reverses a string
def reverse(s):
    if len(s) == 1:
        return s
    else:
        return reverse(s[1:])+s[0]
```

```
>>> print reverse("hello world")
"dlrow olleh"
```

Search a sorted list

(think phonebook)

- How would you search a sorted list to check whether a certain item appears in the list and where?
 - Random search (*yep, I know, this is a stupid algorithm*)
 - Serial search – $O(n)$
 - Binary search



hunting a lion in the desert

Binary search

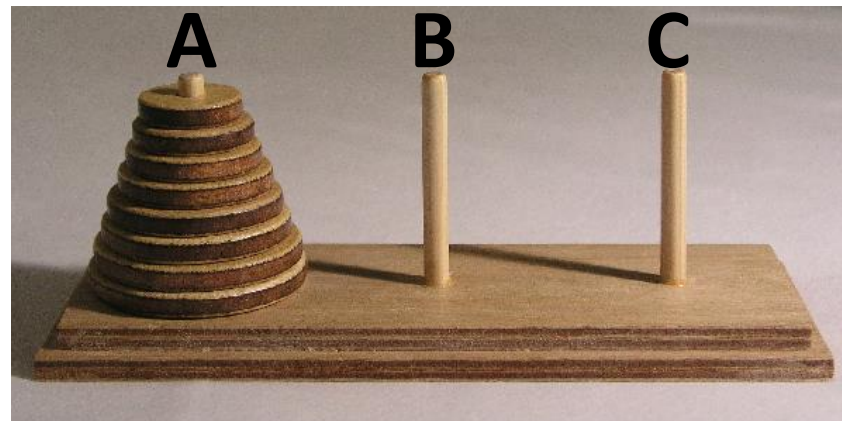
The binary-search algorithm

1. If your list is of size 0, return “not-found”.
2. Check the item located in the middle of your list.
3. If this item is **equal** to the item you are looking for: **you’re done!** Return “found”.
4. If this item is **bigger** than the item you are looking for: do a binary-search on the **first half** of the list.
5. If this item is **smaller** than the item you are looking for: do a binary-search on the **second half** of the list.

How long does it take for this algorithm to find the query item (or to determine it is not in the list)?

Towers of Hanoi

- There are three posts and 64 concentric disks shaped like a pyramid.
- The goal is to move the disks from post **A** to post **B**, following these three rules:
 1. You can move only one disk at a time.
 2. A disk may not be “set aside”. It may only be stacked on one of the three posts.
 3. A larger disk may never be placed on top of a smaller one.



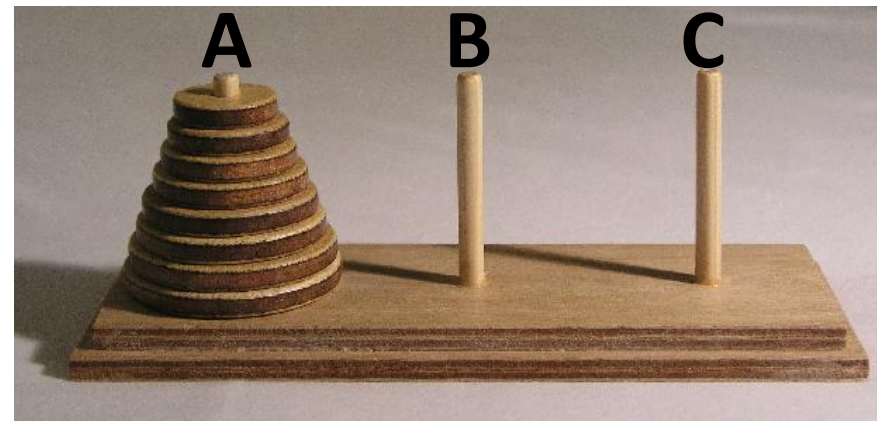
Towers of Hanoi

Towers-of-Hanoi algorithm (for an “n disk tower”)

- 1. Move an “n-1 disk tower” from source-post to resting-post (use the tower-of-hanoi algorithm)*
- 2. Move 1 disk from source-post to destination-post*
- 3. Move an “n-1 disk tower” from resting-post to destination-post (use the tower-of-hanoi algorithm)*

- **What should the base case be?**

- Assuming each disk move takes 1 second, how long would it take to move a 64 disk tower?



**Finally,
let's get back to our merge sort**

The merge sort algorithm

1. Split your list into two halves
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The merge sort algorithm

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4 helper function

```
# Merge two sorted lists
def merge(list1, list2):
    merged_list = []
    i1 = 0
    i2 = 0

    # Merge
    while i1 < len(list1) and i2 < len(list2):
        if list1[i1] <= list2[i2]:
            merged_list.append(list1[i1])
            i1 += 1
        else:
            merged_list.append(list2[i2])
            i2 += 1

    # One list is done, move what's left
    while i1 < len(list1):
        merged_list.append(list1[i1])
        i1 += 1
    while i2 < len(list2):
        merged_list.append(list2[i2])
        i2 += 1

    return merged_list

# merge sort recursive
def sort_r(list):
    if len(list) > 1: # Still need to sort
        half_point = len(list)/2
        first_half = list[:half_point]
        second_half = list[half_point:]

        first_half_sorted = sort_r(first_half)
        second_half_sorted = sort_r(second_half)

        sorted_list = merge \
            (first_half_sorted, second_half_sorted)
        return sorted_list
    else:
        return list
```

List of size 1.
Base case

Recursion vs. Iteration

- There are usually similarities between an iterative solutions (e.g., looping) and a recursive solution.
 - In fact, anything that can be done with a loop can be done with a simple recursive function!
 - In many cases, a recursive solution can be easily converted into an iterative solution using a loop (but not always).
- Recursion can be very costly!
 - Calling a function entails overhead
 - Overhead can be high when function calls are numerous (stack overflow)

Recursion - the take home message

- **Recursion is a great tool to have in your problem-solving toolbox.**
- In many cases, recursion provides a natural and elegant solution to complex problems.
- If the recursive version and the loop version are similar, prefer the loop version to avoid overhead.
- Yet, even in these cases, recursion offers a creative way to **think** about how a problem could be solved.

Sample problem #1

- Write a function that calculates the sum of the elements in a list using a recursion

Hint: your code should not include ANY for-loop or while-loop!

- Put your function in a module, import it into another code file and use it to sum the elements of some list.

Solution #1

utils.py

```
def sum_recursive(a_list):  
    if len(a_list) == 1:  
        return a_list[0]  
    else:  
        return a_list[0] + sum_recursive(a_list[1:])
```

my_prog.py

```
my_list = [1, 3, 5, 7, 9, 11]  
  
from utils import sum_recursive  
print sum_recursive(my_list)
```

Sample problem #2

- Write a **recursive** function that determines whether a string is a palindrome. Again, make sure your code does not include any loops.

A palindrome is a word or a sequence that can be read the same way in either direction.

For example:

- “detartrated”
- “olson in oslo”
- “step on no pets”

Solution #2

```
def is_palindrome(word):  
    l = len(word)  
    if l <= 1:  
        return True  
    else:  
        return word[0] == word[l-1] and is_palindrome(word[1:l-1])
```

```
>>>is_palindrome("step on no pets")  
True  
>>>is_palindrome("step on no dogs")  
False  
>>>is_palindrome("12345678987654321")  
True  
>>>is_palindrome("1234")  
False
```

Challenge problems

1. Write a recursive function that prime factorize s an integer number.

(The prime factors of an integer are the prime numbers that divide the integer exactly, without leaving a remainder).

Your function should print the list of prime factors:

```
>>> prime_factorize(5624)
2 2 2 19 37
>>> prime_factorize(277147332)
2 2 3 3 3 3 3 7 7 11 23 23
```

Note: you can use a for loop to find a divisor of a number but the factorization process itself should be recursive!

2. Improve your function so that it “returns” a list containing the prime factors. Use pass-by-reference to return the list.

3. Can you do it without using ANY loops whatsoever?

Challenge solution 1

```
import math

def prime_factorize(number):

    # find the first divisor
    divisor = number
    for i in range(2, int(math.sqrt(number)) + 1):
        if number % i == 0:
            divisor = i
            break

    print divisor,

    if divisor == number: # number is prime. nothing more to do
        return
    else:                 # We found another divisor, continue
        prime_factorize(number/divisor)

prime_factorize(277147332)
```

Challenge solution 2

```
import math

def prime_factorize(number, factors=[]):

    # find the first divisor
    divisor = number
    for i in range(2, int(math.sqrt(number)) + 1):
        if number % i == 0:
            divisor = i
            break

    factors.append(divisor)

    if divisor == number: # number is prime. nothing more to do
        return
    else:                 # We found another divisor, continue
        prime_factorize(number/divisor, factors)

factors = []
prime_factorize(277147332, factors)
print factors
```

